

Compilers

The graduate version - Spring 2020

Goals

- To become knowledgeable of the foundational concepts underlying modern compiler optimization
- To explore and understand the tradeoffs required when implementing scalable program analyses
- To become familiar with a production-quality compiler system that you can use in your own research

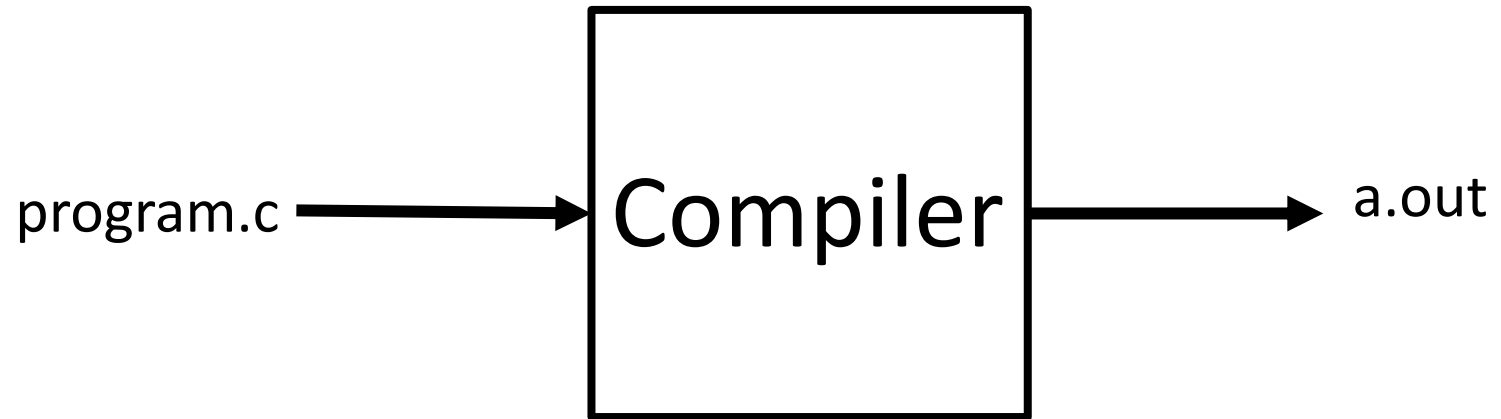
A bit about me ...

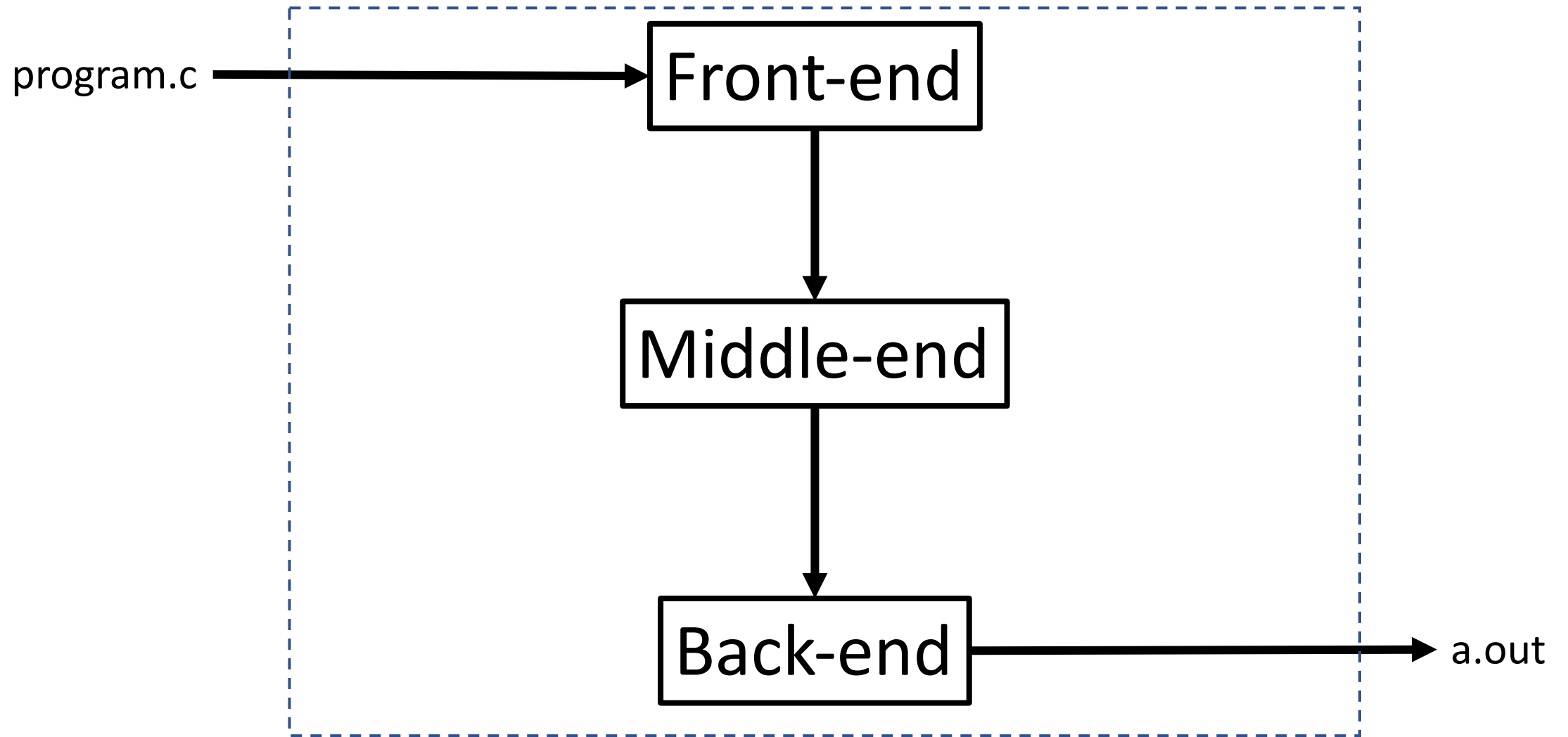
- Worked as a compiler developer in industry from 1986-1990
- Doctoral work on data flow analysis
- Have taken three courses in compilers (all grad courses)
- Have taught undergrad and graduate compilers 20 times
 - 5 different instantiations of the course
- Have implemented significant parts of 7 compilers
 - Most recently this summer (as you will see)
- Lead research on topics that are closely related to compilation

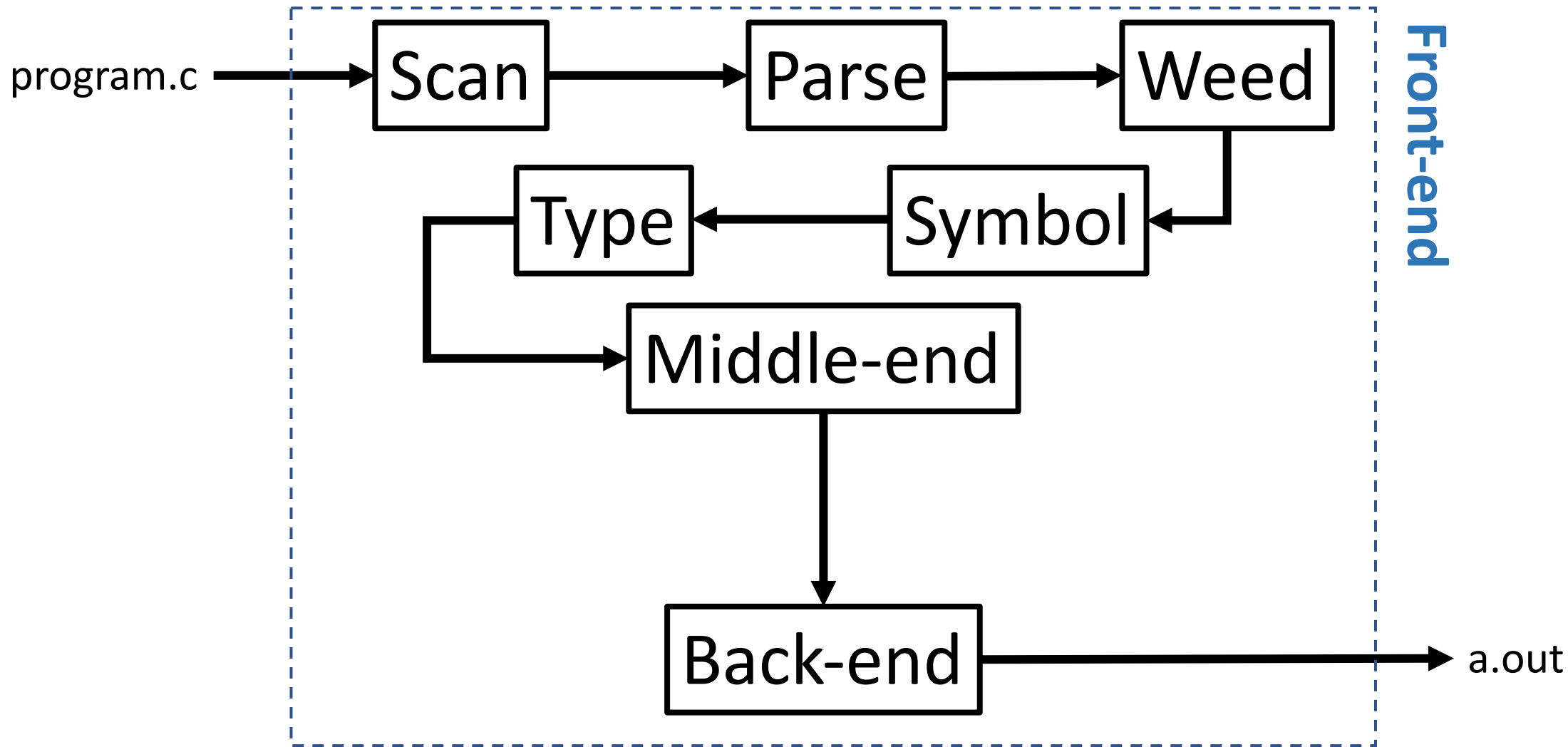
From theory to normal engineering

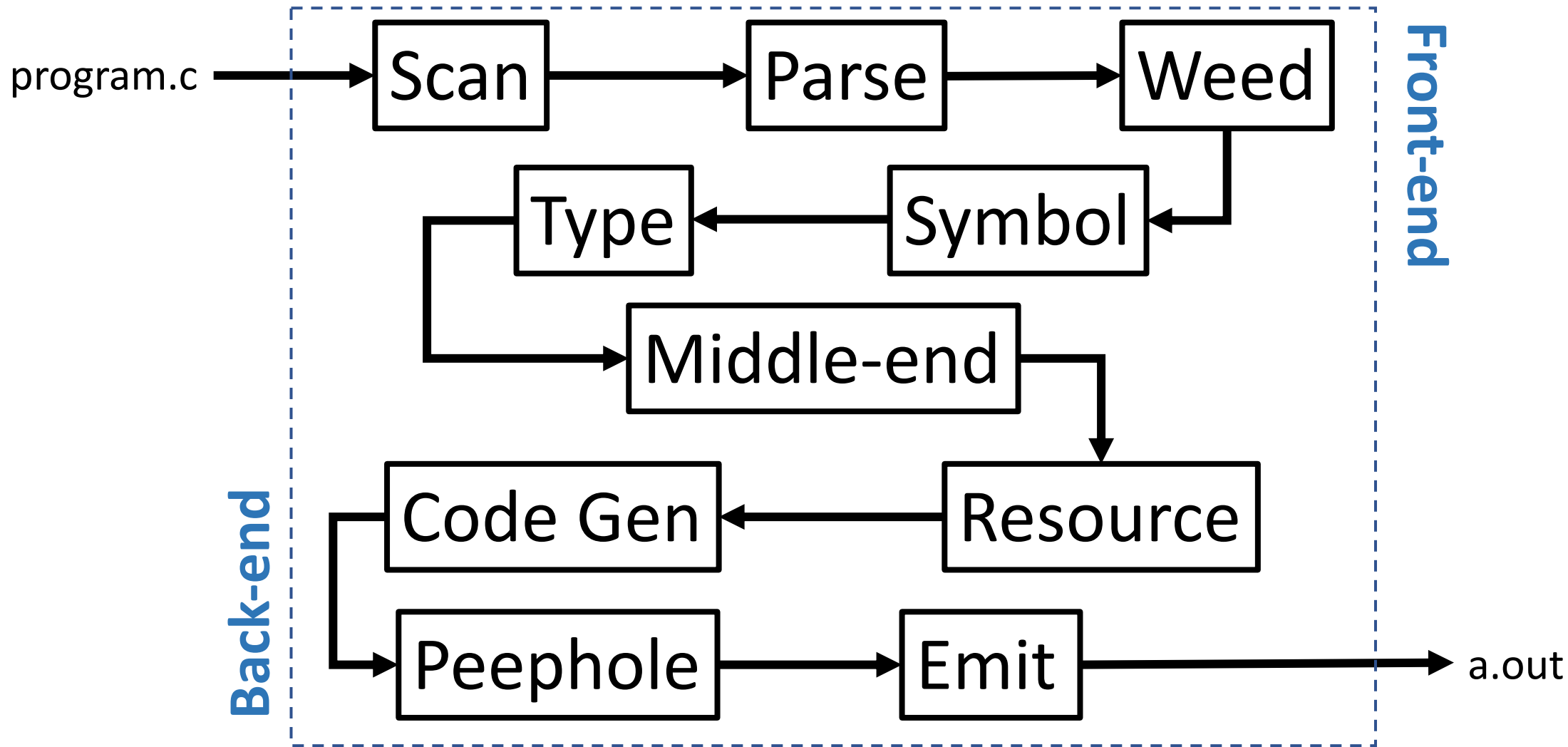
- in the 1960s compilation was *art*
- in the 1970s compilation was *theory*, i.e., studied by theoreticians
- in the 1980s and 90s compilation was *engineering*, i.e., studied as a software product line, supported by reusable programming frameworks and DSLs
- in the 2000s those frameworks became more powerful
- in the 2010s we finally figured out how to test them
- it is one of the most mature software domains you will ever encounter

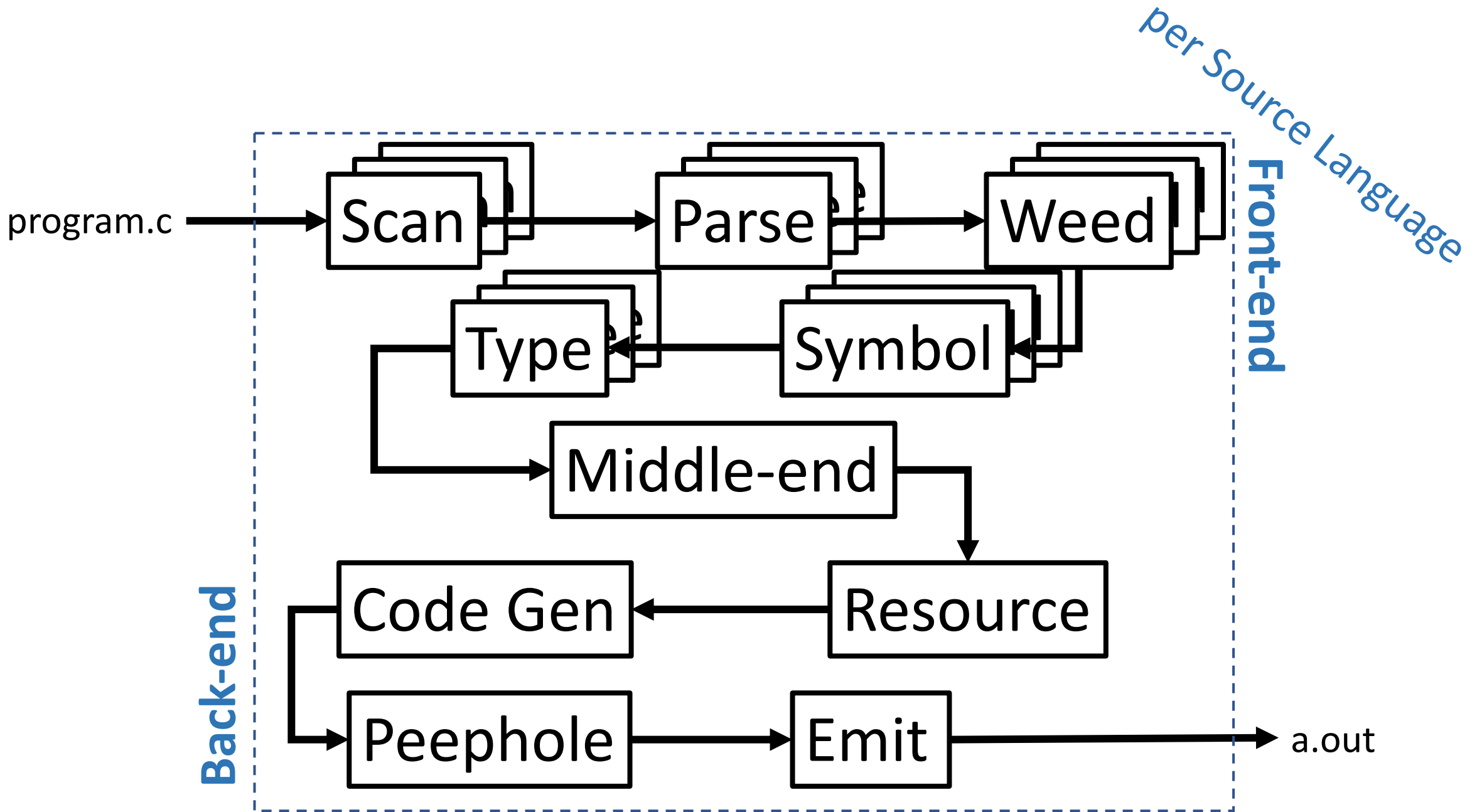
What is a compiler?

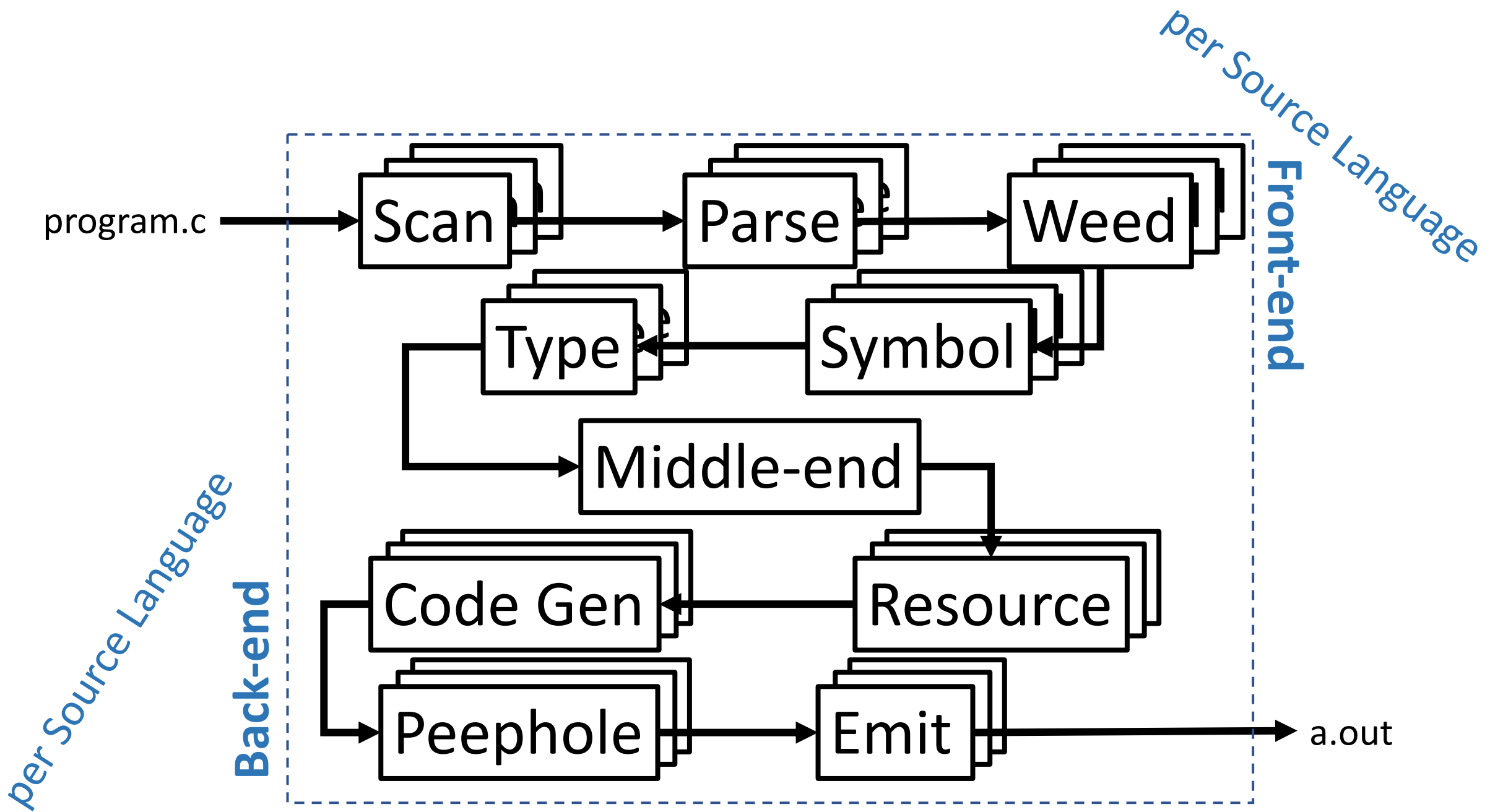






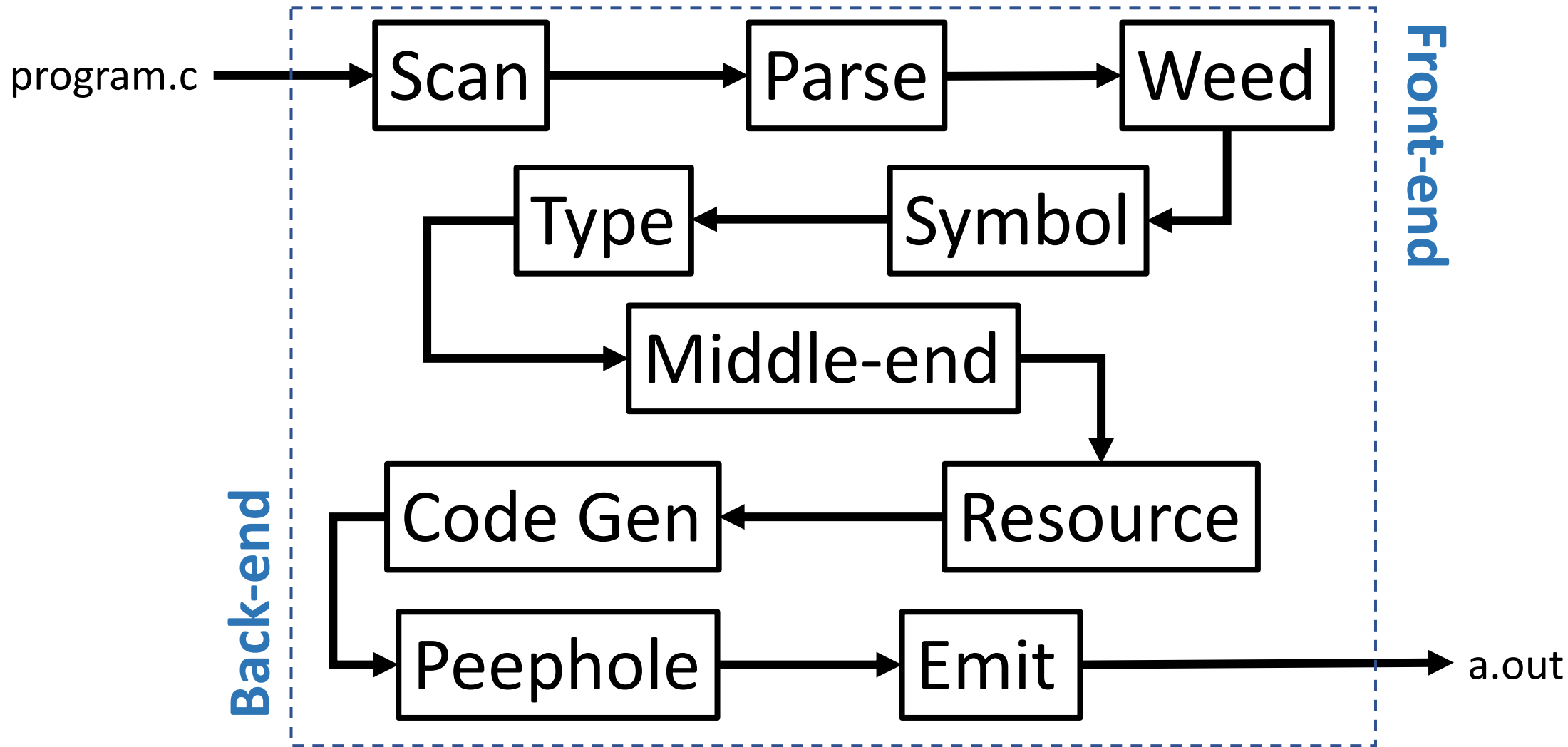




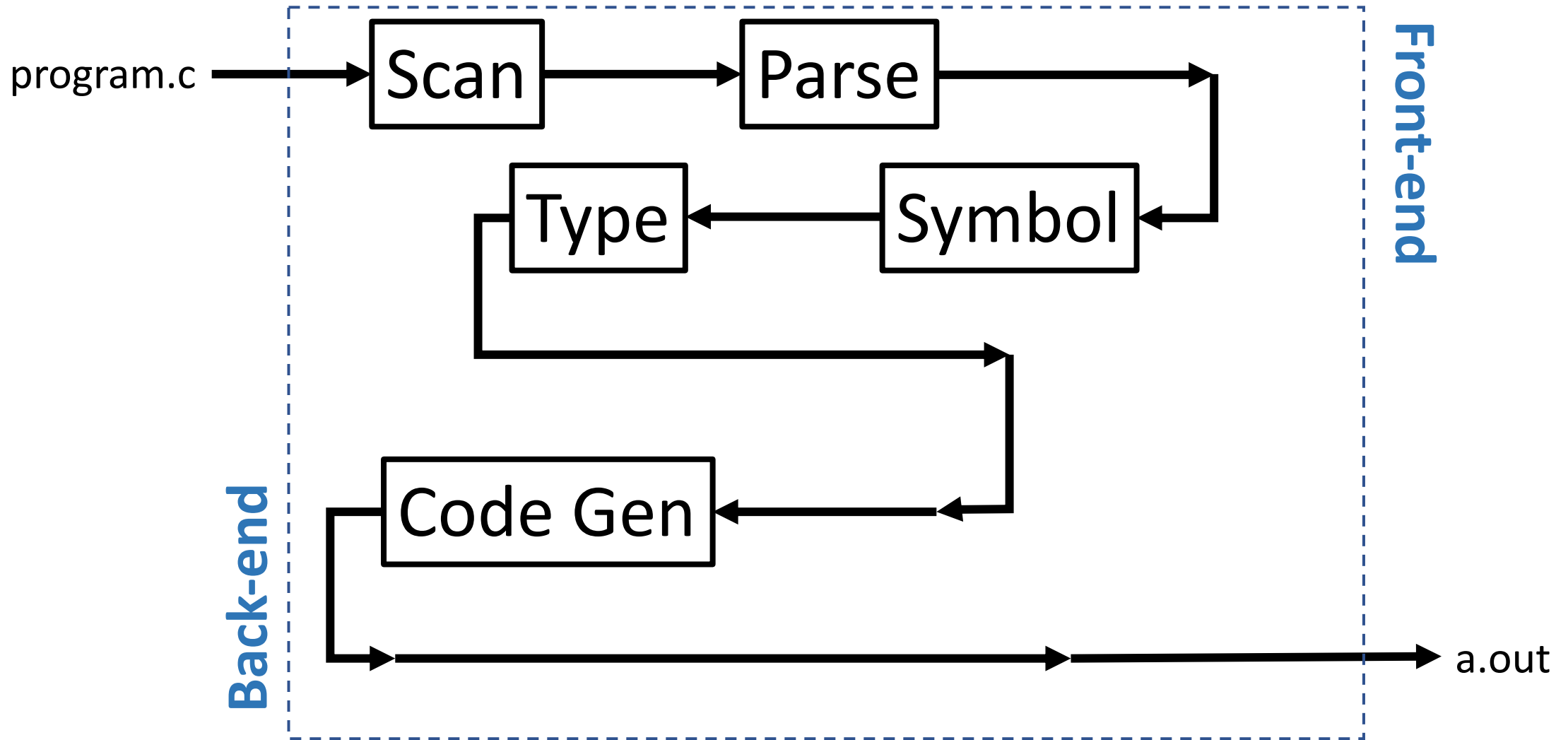


Compilers are ...

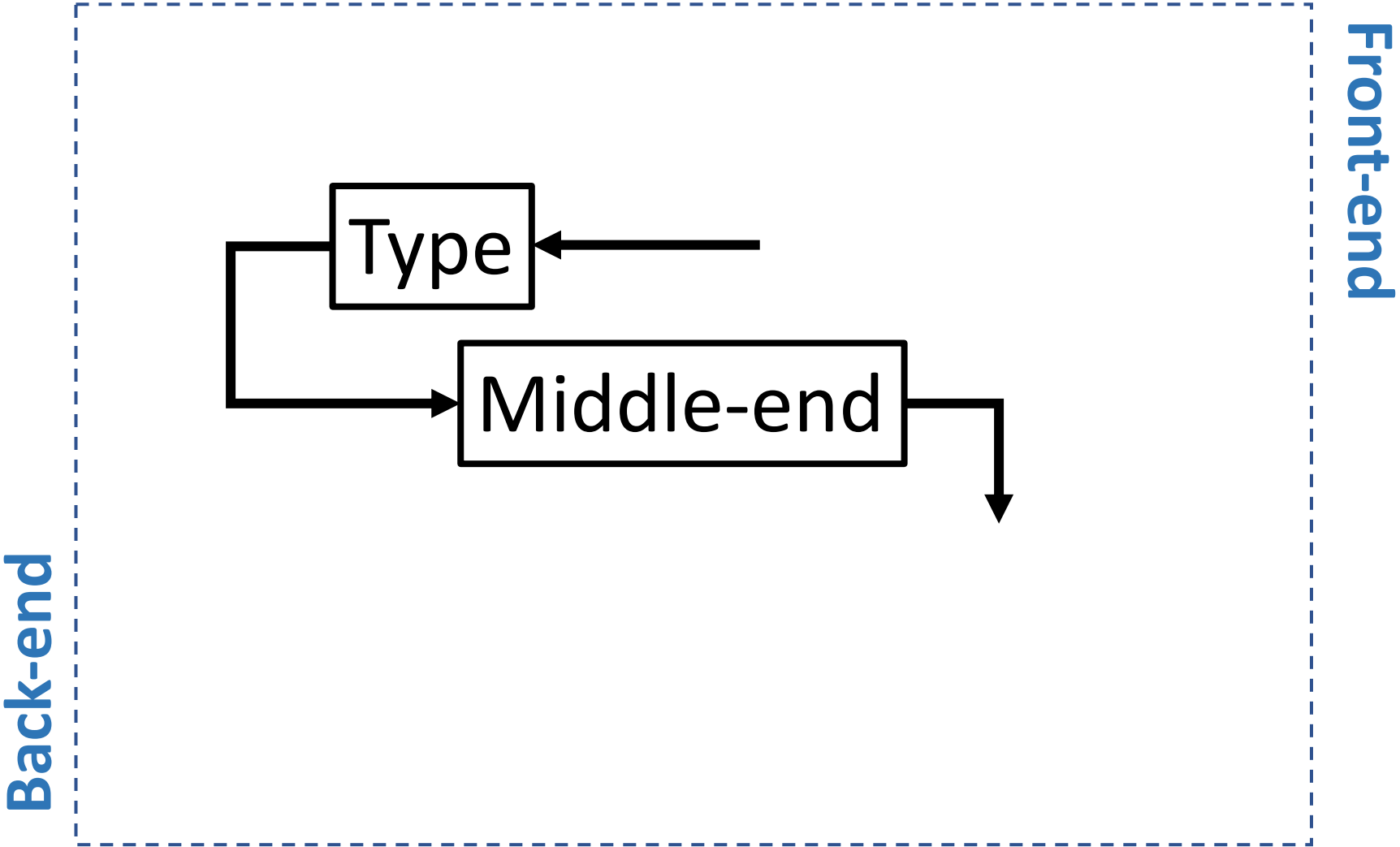
- Large complex software systems
 - GCC >7MSLOC
 - CLANG+LLVM >4MSLOC
- Highly-structured software architectures
 - Well-defined interfaces
 - Components modularized and plug compatible
- Focused on the input and output languages, e.g., for GCC
 - C, C++, Objective C, Ada, Fortran, Go, D, Cobol, Modula-2/3, ...
 - arm, alpha, i386, mips, rs6000, sparc, ... (51 currently)
- **We are going to side-step a lot of that complexity**



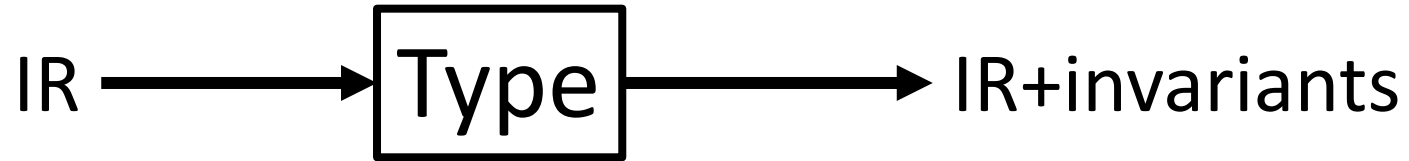
Undergraduate Compilers



This Class



Static Program Analysis



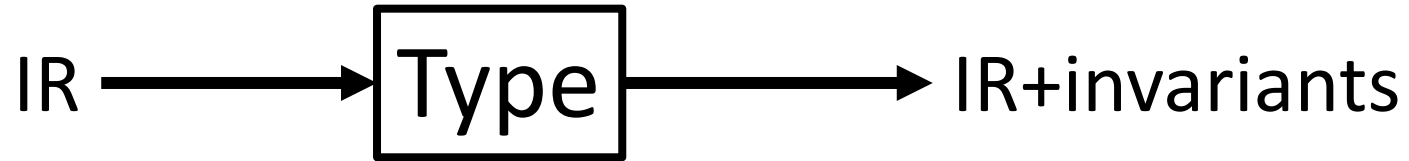
Intermediate Representation

Facts about program behavior that always hold



Static Program Analysis

abstract syntax tree,
symbol table, ...



x is an “integer”
foo(x) returns an “integer”

control flow graph,
dependence graph,
call graph, ...



x+y is always z-10
p and q never point to the same memory
foo() is always called with positive args

Compilers in three parts

- Theory in a controlled environment
 - TIP – Tiny Imperative Language
 - Scala implementation of interpreter and analyses (with holdbacks)
- Practice in the wild
 - `tipc` a compiler from (a subset of) TIP to LLVM bytecode
 - Yours to extend in a class project
- Prompts to drive your exploration and learning
 - Analysis passes in LLVM

A degree of independence will be required

- Theory in a controlled environment
 - TIP is 4500 SLOC of Scala
 - Much of it you will not need to touch or even look at
 - 46 lines marked “`??? //<---` Complete here”
- Practice in the wild
 - `tipc` is about 1000 SLOC of C++
 - Makes heavy use of LLVM APIs and coding idioms (smart pointers)
 - Uses ANTLR4 grammar and custom visitors for AST construction and code-gen
- There is no TA
 - I can be of help for many issues (I implemented `tipc`)
 - I don't use IDEs, so I can't help with that, but I hear they are great